## **Worksheet 1 Solutions:**

(1) Diagram (a) does not represent a finite state machine. In State A, there is a transition on 'a' to both State B and State C. Therefore, the transition function is not a function!

Diagram (b) does describe a finite state machine with transition table on the right. Because we can write a transition table, the transition function  $\delta$  is indeed a function. We can check the other components. The set of state is finite and equals  $\{A,B,C\}$ . There is one starting state  $A.$  The set of accepting states is  $\{C\}.$ 



(2) The input alphabet is  $\{a, c, g, t\}$ . The transition table is on the right.

 $\hat{\delta}(C, ccat) = B$  by chasing arrows. We also chase arrows to determine whether strings are accepted. If we process  $a$   $at$ t $t$ g $t$ t $g$ , we start in A and then go to C (on the second a), to B on t, to B on t, to B on t, to A on g, to C on t, to B on t, to g on A, which is the accepting state.

Processing  $a$  at t t t t gets us from A to B to C to B to B to B to B,  $\,$ which is **not** an accepting state.



(3) We have

$$
\hat{\delta}(C,110110) = \delta(\hat{\delta}(C,11011),0)
$$
\n
$$
= \delta(\delta(\hat{\delta}(C,1101),1),0)
$$
\n
$$
= \delta(\delta(\delta(\hat{\delta}(C,110),1),1),0)
$$
\n
$$
= \delta(\delta(\delta(\delta(\hat{\delta}(C,11),0),1),1),0)
$$
\n
$$
= \delta(\delta(\delta(\delta(\delta(\hat{\delta}(C,1),1),0),1),1),0)
$$
\n
$$
= \delta(\delta(\delta(\delta(\delta(C,1),1),0),1),1),0)
$$
\n
$$
= \delta(\delta(\delta(\delta(\delta(D,1),0),1),1),0)
$$
\n
$$
= \delta(\delta(\delta(\delta(D,1),0),1),1),0)
$$
\n
$$
= \delta(\delta(\delta(C,1),1),0)
$$
\n
$$
= \delta(\delta(D,1),0)
$$
\n
$$
= \delta(E,0)
$$
\n
$$
= C.
$$

(4) The input alphabet of the new deterministic automaton is  $\{0,1\}.$  If we start in A, we only have one transition on 0 to B, but then an  $\epsilon$ -transition from B to C (which we can but do not have to take), therefore in the new deterministic automaton:

$$
\delta({A},0) = {B}, C.
$$

 $\delta({A},1) = { }$ .

Similarly, if we start in B, we can take an  $\epsilon$ -transition before or after another transition. Therefore:

 $\delta({B},0) = {B,C}$ 

and

 $\delta({B},1) = {D}.$ 

Finally, there is no transition out of C on 0, and only one possibility to transition on 1, so that  $\delta({C}, 0) = {},$ 

 $\delta({C},1) = {D}.$