

Finite State Machines and String Matching

String Matching

- Idea:
 - We look at each character in the pattern exactly once
 - We use a finite state machine to “remember” what we have matched
 - We create the finite state machine by pre-processing the pattern

String Matching

- Example:
- Pattern is

A	C	A	G	A	A	T
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- States corresponds to prefixes of the string

Seen
""

Seen
"A"

Seen
"AC"

Seen
"ACA"

Seen
ACAG

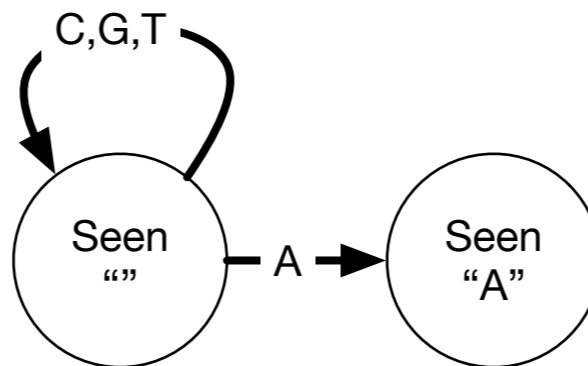
Seen
"ACA
GA"

Seen
"ACA
GAA"

Seen
"ACA
GAAT"

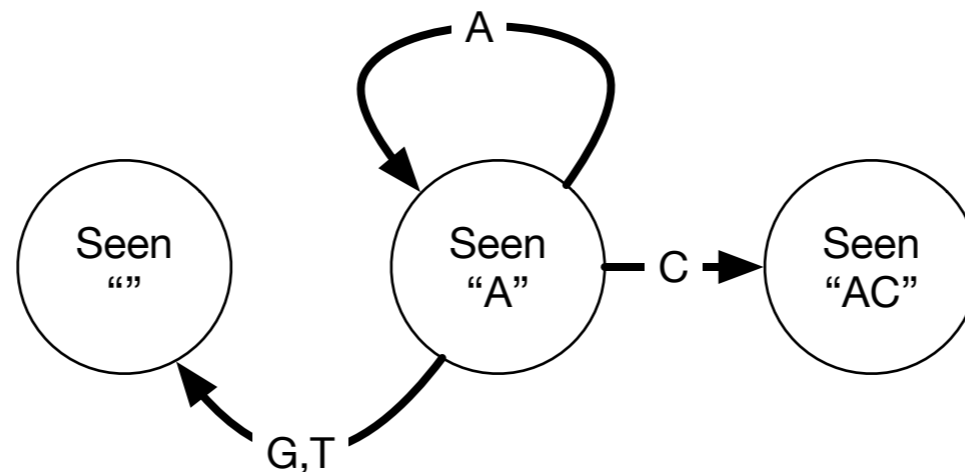
String Matching

- State 0 corresponds to nothing currently matched
 - When we match an “A”, we move to the next State
 - Otherwise, we go back to the current state



String Matching

- In State “Seen A”:
 - If we have a “C”, go to next State
 - If we have a G or T, go back
 - If we have an “A”, we have seen an “AA” which has a suffix that partially matches



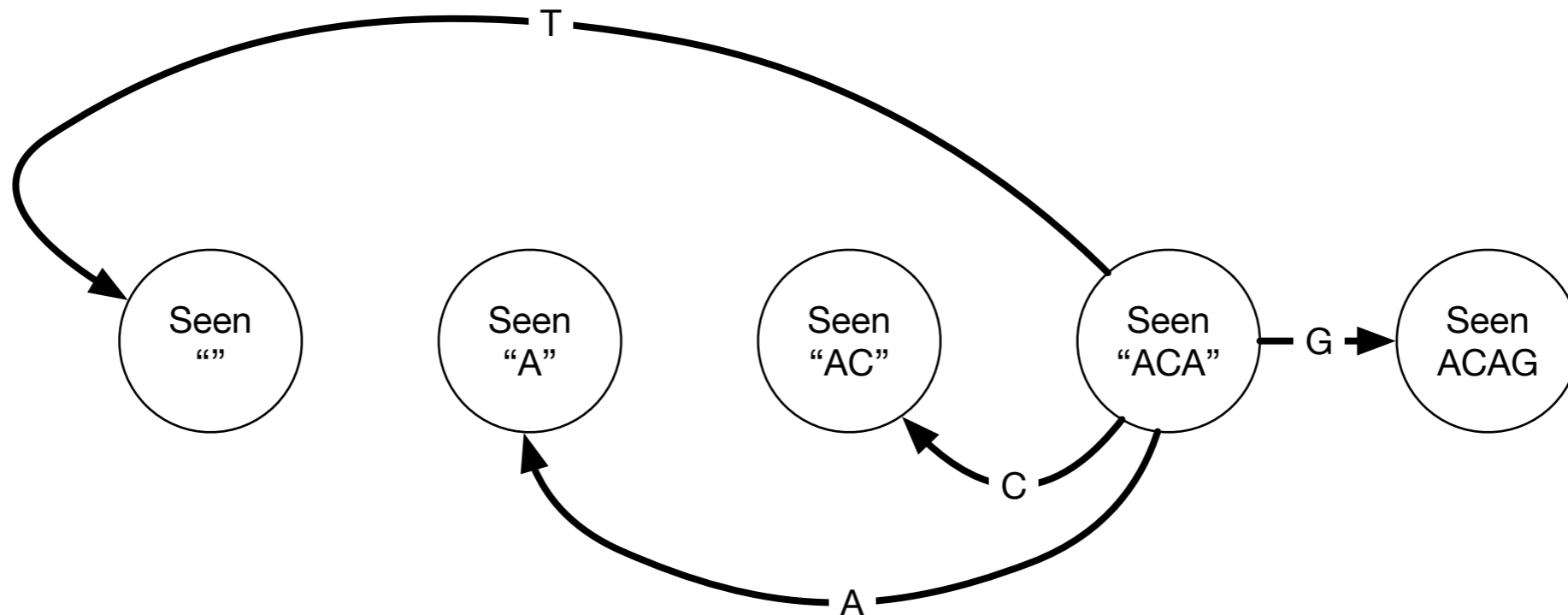
String Matching

- State “AC”
 - “A” move to “ACA”
 - “C”, “G”, “T” move to “ ”



String Matching

- State “ACA”
 - “C” matches “AC”
 - “T” matches “”
- State “ACA”
 - “G” matches next state
 - “A” matches “A”



String Matching

- Your turn: Quiz