Data at Scale

Remote Storage

- High Availability Measures
	- Replication
	- Erasure Correcting Coding

• Data broken into pages

$$
P_1 =
$$

$$
P_2 =
$$

$$
P_3 =
$$

.

. . .

 $\{p_{1,0}, p_{1,1}, p_{1,2}, \ldots\}$ $\{p_{2,0}, p_{2,1}, p_{2,2}, \ldots\}$ $\{p_{3,0}, p_{3,1}, p_{3,2}, \ldots\}$

• Create parity pages . .

$$
C_1 = \{c_{1,0}, c_{1,1}, c_{1,2}, \ldots\}
$$

\n
$$
C_2 = \{c_{2,0}, c_{2,1}, c_{2,2}, \ldots\}
$$

\n
$$
C_3 = \{c_{3,0}, c_{3,1}, c_{3,2}, \ldots\}
$$

• "Check" pages calculated using an erasure correcting code

$$
(c_{1,i},c_{2,i},\ldots,c_{m,i})=\Phi(p_{1,i},p_{2,i},\ldots,p_{n,i})
$$

- *m*/*n-*code:
	- *• m* data pages
	- *• n* check pages
	- *•* Any *m* of the *n*+*m* pages are sufficient to reconstruct all *m* data pages

- Related to error-correcting codes in telecommunications / networking
- Simplest code: Parity code
	- One check symbol
	- Is parity of the data symbols

$$
c_{1,i}=p_{1,i}\oplus p_{2,i}\oplus\ldots\oplus p_{n,i}
$$

• Can calculate single lost data symbol as the parity of the survivors and the check symbol

- Group Activity:
	- Reconstruct the missing text

- Many codes are known
- Some involve algebraic objects such as finite fields
	- Example GF(2^{^8})
		- Elements are bytes (string of 8 bits)
		- Addition is exclusive-or
			- 0010 0011 + 1011 1110 = 1001 1101
		- Multiplication can be defined as polynomial multiplication $00100011 \approx t^5 + t + 1$

- Multiplication in GF(2^{^8}) $00100011 \approx t^5 + t + 1$ $10111110 \approx t^7 + t^5 + t^4 + t^3 + t^2 + t$
	- Multiply polynomials with coefficients in $\{0,1\}$
		- Caution: $1+1=0$

 $t^{12} + t^{10} + t^{9} + t^{8} + t^{7} + t^{6}$ $+t^8+t^6+t^5+t^4+t^3+t^2$ $+t^7+t^5+t^4+t^3+t^2+t^2$ $= t^{12} + t^{10} + t^9 + t^8$

• Now divide this polynomial by a *generator polynomial*

$$
t^8 + t^4 + t^3 + t + 1
$$

- The remainder is the product
	- Group exercise: Calculate the remainder

- However, we do not have to do this multiplication every time
	- Can use tables and algebraic properties
	- Plank, Greenan, Miller: Can multiply 16 GF elements by a constant GF element with 6 assembly instructions
		- Using PSHUFB instruction

- Simple method to create erasure correcting codes
	- Start with Vandermonde matrix of size *n* by *n*+*m*
		- Defined by $n+m$ different Galois field elements

$$
\mathbf{V} = \begin{pmatrix} 1 & 1 & 1 & 1 & \dots & 1 \\ \alpha_0 & \alpha_2 & \alpha_3 & \alpha_4 & \dots & \alpha_{n+m-1} \\ \alpha_0^2 & \alpha_2^2 & \alpha_3^2 & \alpha_4^2 & \dots & \alpha_{n+m-1}^2 \\ \alpha_0^3 & \alpha_2^3 & \alpha_3^3 & \alpha_4^3 & \dots & \alpha_{n+m-1}^3 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ \alpha_0^{n-1} & \alpha_2^{n-1} & \alpha_3^{n-1} & \alpha_4^{n-1} & \dots & \alpha_{n+m-1}^{n-1} \end{pmatrix}
$$

• Use elementary row transformations to obtain

$$
\mathbf{G} = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 & r_{0,0} & r_{0,1} & \dots & r_{0,m} \\ 0 & 1 & 0 & \dots & 0 & r_{1,0} & r_{1,1} & \dots & r_{1,m} \\ 0 & 0 & 1 & \dots & 0 & r_{2,0} & r_{2,1} & \dots & r_{2,m} \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 & r_{n,0} & r_{n,1} & \dots & r_{n,m} \end{pmatrix}
$$

• Always possible

- This matrix has the remarkable property that any *n* by *n* sub-matrix is invertible
- Follows that multiplication with **G** defines an erasure correcting code
	- **• d** *n* data symbols written as a row vector
	- **• c** the same *n* symbols followed by *m* check symbols

$$
\mathbf{c} = \mathbf{d} \cdot \mathbf{G}
$$

- These "linear" codes have remarkable properties
	- Can change a single page

•

• Can calculate the new check pages from the old and the new values of the page

- Are used in RAID Level 6 storage systems
- Can / could be used across system boundaries