### No SQL Databases

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#### Relational Model Shortcomings

- Greater Scalability
	- High write throughput / very large datasets
- Independence from few vendors Move towards Open **Source**
- Need for different query operations
- Restrictiveness of relational schemas

- Hush: HBase URL Shortener
	- Hand a URL to a Shortener service
	- Get a shorter URL back
		- E.g. to use in twitter messages
	- Shortener provides counter for each shortened URLs
	- "Vanity URL" that incorporate specific domain names
	- Need to maintain users
		- log in to create short URLs
		- track existing URLs
		- see reports for daily, weekly, or monthly usage

- Data is too large to store at a single server
	- Notice:
		- Limited need for transactions
		- Importance of high throughput writes and reads

- Columnar Layout
	- A relational database strategy often adopted in No-SQL databases
	- Instead of storing data in tuples
	- Store by attribute

- For large HUSH:
	- Can use a relational database
	- Use normalization and obtain a scheme

- Principles of Denormalization, Duplication, Intelligent Keys
	- Denormalize by duplicating data in more than one table
		- Avoids aggregation at read time
		- Prematerialize required views

- Example: HBase URL Shortener (Hush)
	- user(id, username, credentials, rules, first\_name, last\_name, email) with unique username constraint
	- url(id, url, refShortID, title, description, content)
	- shorturl(id, userID, urIID, shortID, refShortID, description) with unique shortID and F.K. userID and urlID
	- click(id, datestamp, shortID, category, dimension, counter) with F.K. shortID

• Purpose: maps long URLs to short URLs



- Short URL can be given to others
- This is translated to the full URL



• Each click is tracked, which aggregates to weekly usage numbers



• All these operations require joins



- Bandwidth problem
	- Especially for joins
	- Need to store data in joins together, not look them up separately
	- But can relax on the consistency model:
		- No need to serialize short URL creation or URL translations or have atomic updates
	- Might be able to relax integrity constraints
		- Statistics need to be approximately correct

- Denormalization:
	- Key idea: Store data together that is likely to be joined
	- Means:
		- massive duplication of data
		- relaxed consistency needed
		- but faster reads / writes

- Wide column stores
	- names and format of columns can vary from row to row
	- Google's BigTable

- Document databases
	- MongoDB, XML databases
		- see below

- Key-value database
	- Every record is a key-value pair
	- A large set of tools predating no-sql databases in general

- Graph databases
	- Navigational database successor:
		- information about data interconnectivity or topology as important as data itself
	- See below for an example

• Data is often structured hierarchically

```
Invoice = \{date : "2008-05-24"
   invoiceNumber : 421
   InvoiceItems : {
     Item : {
       description : "Wool Paddock Shet Ret Double Bound Yellow 4'0"
       quantity : 1
       unitPrice : 105.00
     }
     Item : {
       description : "Wool Race Roller and Breastplate Red Double"
       quantity : 1
       unitPrice : 75.00
     }
     Item : {
       description : "Paddock Jacket Red Size Medium Inc Embroidery"
       quantity : 2
       unitPrice : 67.50
   }
   }
}
```
#### • As an XML document

<invoice>

```
<number>421</number>
\text{Value}>2008-05-24 </date>
 <items>
   <item>
    <description>Wool Paddock Shet Ret Double Bound Yellow 4'0"</description>
    <quantity>1</quantity>
    <unitPrice>105.00</unitPrice>
  \langleitem\rangle <item>
    <description>Wool Race Roller and Breastplate Red Double</description>
    <quantity>1</quantity>
    <unitPrice>75.00</unitPrice>
  \langleitem\rangle <item>
    <description>Paddock Jacket Red Size Medium Inc Embroidery</description>
    <quantity>2</quantity>
    <unitPrice>67.50</unitPrice>
   </item>
</items>
</invoice>
```
- Advantage of XML
	- Faster to scan all data
	- No joins
- Disadvantages of XML
	- Each record contains the full or an abbreviated scheme
	- Each query needs to select from big chunks of data

- JSON JavaScript Object Notation
	- Human-readable
	- Organized as key-value pairs

}

• JSON record example

```
 "firstName": "John",
   "lastName": "Smith",
   "isAlive": true,
   "age": 27,
   "address": {
     "streetAddress": "21 2nd Street",
     "city": "New York",
     "state": "NY",
     "postalCode": "10021-3100"
   },
   "phoneNumbers": [
 {
       "type": "home",
       "number": "212 555-1234"
     },
     {
       "type": "office",
       "number": "646 555-4567"
     },
     {
       "type": "mobile",
       "number": "123 456-7890"
     }
  \mathbf{1},
   "children": [],
   "spouse": null
```
- JSON can use a schema (type definition)
- JSON was first used for data transmission as a data serialization format

- Many-to-One and Many-to-Many Relationships
	- Modeled by the same value for the same key
		- Problem: Need to standardize / internationalize these values
		- Using id-s instead of plain text to avoid problems
		- Table of id-s reintroduce a relational scheme through a backdoor

- Resumé
	- Users present people
	- People have jobs, education, and recommenders
- But they share jobs, companies, degrees, schools, recommenders
	- Should they stay text strings or become entities?
		- Latter allows to add information to all resumés
- If recommenders get a photo, then all resumés should be updated with this photo, so better to make recommenders entities

Data has a tendency to become less-join free



- Records are documents
	- Encode in
		- XML
		- YAML
		- JSON
		- BSON (Mongo DB)
	- CRUD operations: create, read, update, delete

- Enforcing schema
	- Most document databases do not enforce schema
		- $\rightarrow$  "Schemaless"
		- In reality: "Schema on Read"
		- RDBMS would then use "Schema on Write"
- Allows schema updates in simple form

- Schema on Read:
	- Advantages:
		- Data might come from external sources
	- Disadvantages:
		- No data checking

- Document database support
	- Most commercial database systems now support XML databases

# Query Languages

- Documents lend themselves to object-oriented querying
	- Imperative code
- SQL is declarative:
	- Programmer explains a solution
	- System figures out the best way to find the solution
- Use declarative query languages for document databases

# Query Languages

- Map-Reduce (neither declarative nor imperative):
	- Consists of only two pieces of code
		- Mapping: Selecting from Documents
		- Reducing: Take selection elements and operate on them

- Graphs consists of vertices and edges
	- Example:
		- Social graphs: vertices are people and edges are relationships such "knows"
		- Web graph: vertices are pages and edges are links
		- Road networks: vertices are places and edges are connections

• Relational Database hides semantic relationships



• Document model hides semantic relationships



- Some property values are really references to foreign aggregates
	- Aggregate's identifier is a foreign key
- Relationships between them are not explicitly accessible
	- Joining aggregates becomes expensive



- Relational Database
	- Some queries are simple:

SELECT p1.Person FROM Person p1 JOIN PersonFriend ON PersonFriend.FriendID = p1.ID JOIN Person p2 ON PersonFriend.PersonID = p2.ID WHERE p2.Person = 'Bob'



- Relational Database
	- Some queries others are more involved: Friends of Bob

SELECT p1.Person FROM Person p1 JOIN PersonFriend

- - ON PersonFriend.PersonID = p1.ID JOIN Person p2
	- ON PersonFriend.FriendID = p2.ID

```
WHERE p2.Person = 'Bob'
```


• Relational Database

• Some queries others are difficult: Alice's friends of friends SELECT p1. Person AS PERSON, p2. Person AS FRIEND OF FRIEND FROM PersonFriend pf1 JOIN Person p1 ON pf1.PersonID = p1.ID JOIN PersonFriend pf2 ON pf2.PersonID = pf1.FriendID JOIN Person p2 ON pf2.FriendID = p2.ID WHERE p1.Person = 'Alice' AND pf2.FriendID <> p1.ID



- Property graph model by Neon
	- Each vertex consists of
		- A unique identifier
		- A set of outgoing edges
		- A set of incoming edges
		- A collection of properties key-value pairs
	- Each edge consists of
		- A unique identifier
		- The tail vertex
		- The head vertex
		- A label to describe the relationship
		- A collection of properties key-value pairs



• Order history as a property graph



- Processing queries in Neo4j
	- Use Cypher (from "The matrix")
	- Can describe a path



 $(emi1) < -$ [:KNOWS]-(jim)-[:KNOWS]->(ian)-[:KNOWS]->(emil)



```
(emil:Person {name:'Emil'}) 
       <-[:KNOWS]-(jim:Person {name:'Jim'}) 
       -[:KNOWS]->(ian:Person {name:'Ian'}) 
      -\lceil: KNOWS] -> (emil)
```
• Finding the mutual friends of Jim:

MATCH  $(a:Person \{name: 'Jim' }) - [ : KNOWS ] -> (b) - [ : KNOWS ] -> (c)$ ,  $(a)$  - $[$ : KNOWS] -> (c) RETURN b, c

• Triple Stores

- Information is stored as (subject, predicate, object)
	- Subjects correspond to vertices
	- Objects are
		- A value in a primitive data type ( jim : age : **64**)
		- Another vertex (jim : friend of : thomas)



- \_:lucy a :Person
- :lucy :name "Lucy"
- :lucy :born in :idaho
- 
- \_:idaho :name "Idaho"
- \_:idaho :type "State"
	-
- 
- 
- 
- \_:idaho :within \_:usa
- 
- 
- 
- :idaho a :Location
	-
	-
	-

- Triple stores are the language of the semantic web
- Semantic web:
	- Machine readable description of type of links
		- e.g. image, text, …
	- Creates web of data a database of everything
- Stored in Resource Description Framework (RDF)
- SPARQL query language for triple stores