#### No SQL Databases

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# Relational Model Shortcomings

- Greater Scalability
  - High write throughput / very large datasets
- Independence from few vendors Move towards Open Source
- Need for different query operations
- Restrictiveness of relational schemas

- Hush: HBase URL Shortener
  - Hand a URL to a Shortener service
  - Get a shorter URL back
    - E.g. to use in twitter messages
  - Shortener provides counter for each shortened URLs
  - "Vanity URL" that incorporate specific domain names
  - Need to maintain users
    - log in to create short URLs
    - track existing URLs
    - see reports for daily, weekly, or monthly usage

- Data is too large to store at a single server
  - Notice:
    - Limited need for transactions
    - Importance of high throughput writes and reads

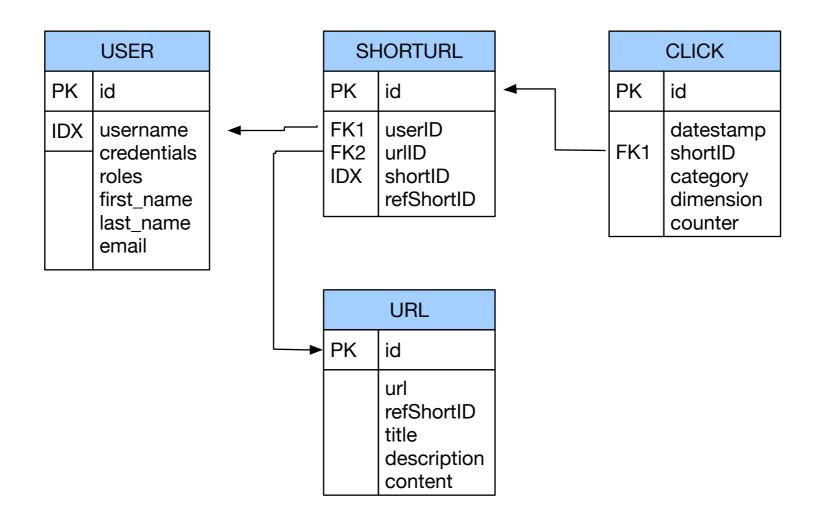
- Columnar Layout
  - A relational database strategy often adopted in No-SQL databases
  - Instead of storing data in tuples
  - Store by attribute

- For large HUSH:
  - Can use a relational database
  - Use normalization and obtain a scheme

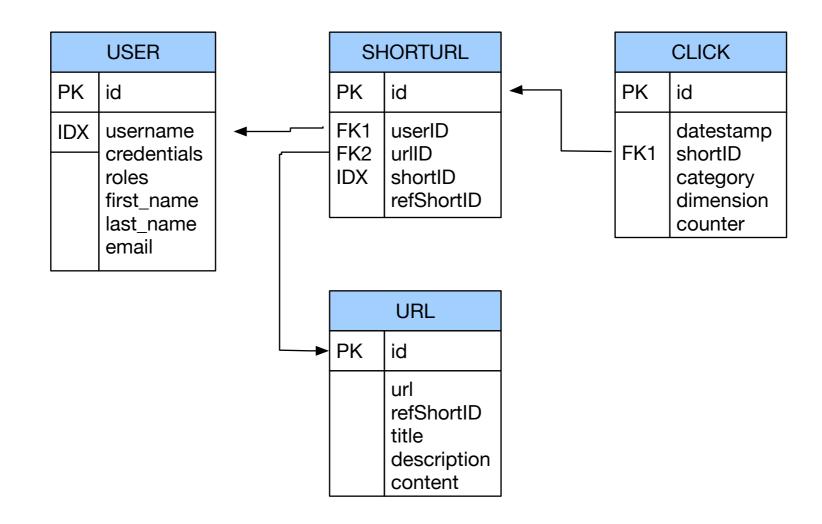
- Principles of Denormalization, Duplication, Intelligent Keys
  - Denormalize by duplicating data in more than one table
    - Avoids aggregation at read time
    - Prematerialize required views

- Example: HBase URL Shortener (Hush)
  - user(<u>id</u>, username, credentials, rules, first\_name, last\_name, email) with unique username constraint
  - url(<u>id</u>, url, refShortID, title, description, content)
  - shorturl(id, userID, urIID, shortID, refShortID, description) with unique shortID and F.K. userID and urIID
  - click(<u>id</u>, datestamp, shortID, category, dimension, counter) with F.K. shortID

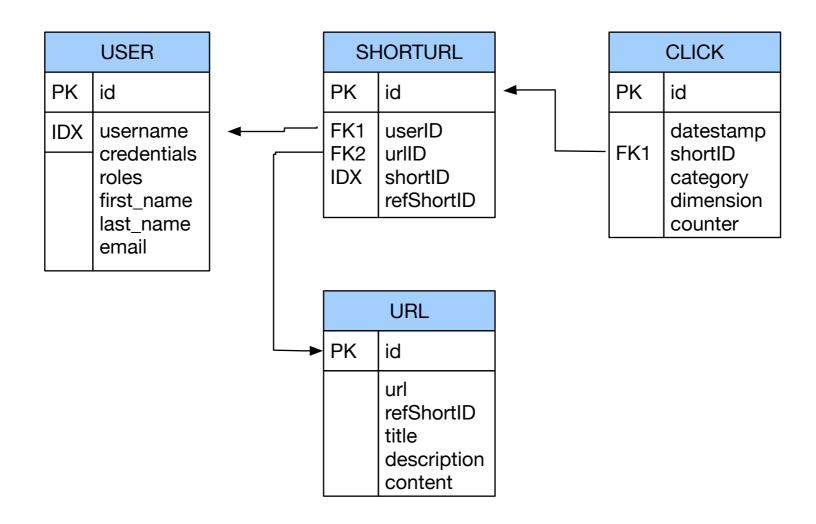
Purpose: maps long URLs to short URLs



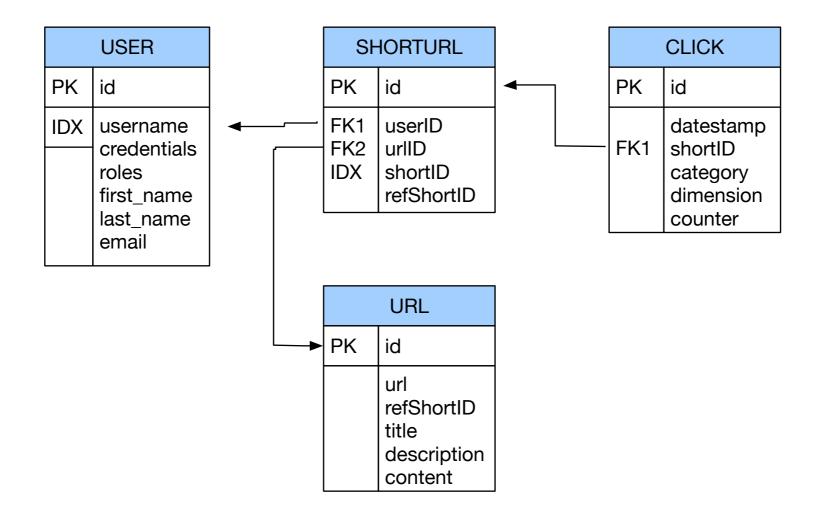
- Short URL can be given to others
- This is translated to the full URL



 Each click is tracked, which aggregates to weekly usage numbers



All these operations require joins



- Bandwidth problem
  - Especially for joins
  - Need to store data in joins together, not look them up separately
  - But can relax on the consistency model:
    - No need to serialize short URL creation or URL translations or have atomic updates
  - Might be able to relax integrity constraints
    - Statistics need to be approximately correct

- Denormalization:
  - Key idea: Store data together that is likely to be joined
  - Means:
    - massive duplication of data
    - relaxed consistency needed
    - but faster reads / writes

- Wide column stores
  - names and format of columns can vary from row to row
  - Google's BigTable

- Document databases
  - MongoDB, XML databases
    - see below

- Key-value database
  - Every record is a key-value pair
  - A large set of tools predating no-sql databases in general

- Graph databases
  - Navigational database successor:
    - information about data interconnectivity or topology as important as data itself
  - See below for an example

Data is often structured hierarchically

```
Invoice = {
 date : "2008-05-24"
  invoiceNumber: 421
  InvoiceItems : {
   Item : {
     description: "Wool Paddock Shet Ret Double Bound Yellow 4'0"
     quantity: 1
     unitPrice: 105.00
    Item : {
     description: "Wool Race Roller and Breastplate Red Double"
     quantity: 1
     unitPrice: 75.00
    Item : {
     description: "Paddock Jacket Red Size Medium Inc Embroidery"
     quantity: 2
     unitPrice: 67.50
```

As an XML document

```
<invoice>
 <number>421</number>
<date>2008-05-24</date>
 <items>
 <item>
  <description>Wool Paddock Shet Ret Double Bound Yellow 4'0"</description>
  <quantity>1</quantity>
  <unitPrice>105.00</unitPrice>
 </item>
 <item>
  <description>Wool Race Roller and Breastplate Red Double/description>
  <quantity>1</quantity>
  <unitPrice>75.00</unitPrice>
 </item>
 <item>
  <description>Paddock Jacket Red Size Medium Inc Embroidery</description>
  <quantity>2</quantity>
  <unitPrice>67.50</unitPrice>
 </item>
 </items>
</invoice>
```

- Advantage of XML
  - Faster to scan all data
  - No joins
- Disadvantages of XML
  - Each record contains the full or an abbreviated scheme
  - Each query needs to select from big chunks of data

- JSON JavaScript Object Notation
  - Human-readable
  - Organized as key-value pairs

JSON record example

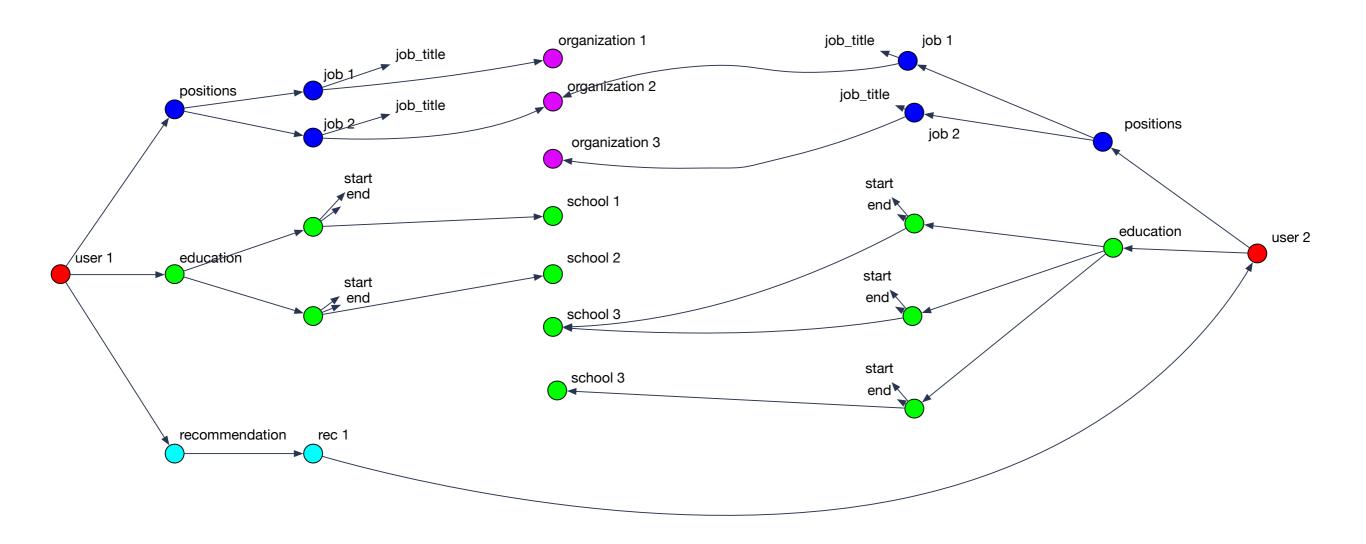
```
"firstName": "John",
"lastName": "Smith",
"isAlive": true,
"age": 27,
"address": {
  "streetAddress": "21 2nd Street",
  "city": "New York",
  "state": "NY",
  "postalCode": "10021-3100"
"phoneNumbers": [
    "type": "home",
    "number": "212 555-1234"
 },
    "type": "office",
    "number": "646 555-4567"
  },
    "type": "mobile",
    "number": "123 456-7890"
"children": [],
"spouse": null
```

- JSON can use a schema (type definition)
- JSON was first used for data transmission as a data serialization format

- Many-to-One and Many-to-Many Relationships
  - Modeled by the same value for the same key
    - Problem: Need to standardize / internationalize these values
    - Using id-s instead of plain text to avoid problems
    - Table of id-s reintroduce a relational scheme through a backdoor

- Resumé
  - Users present people
  - People have jobs, education, and recommenders
- But they share jobs, companies, degrees, schools, recommenders
  - Should they stay text strings or become entities?
    - Latter allows to add information to all resumés
- If recommenders get a photo, then all resumés should be updated with this photo, so better to make recommenders entities

Data has a tendency to become less-join free



- Records are documents
  - Encode in
    - XML
    - YAML
    - JSON
    - BSON (Mongo DB)
  - CRUD operations: create, read, update, delete

- Enforcing schema
  - Most document databases do not enforce schema
    - —> "Schemaless"
    - In reality: "Schema on Read"
    - RDBMS would then use "Schema on Write"
- Allows schema updates in simple form

- Schema on Read:
  - Advantages:
    - Data might come from external sources
  - Disadvantages:
    - No data checking

- Document database support
  - Most commercial database systems now support XML databases

### Query Languages

- Documents lend themselves to object-oriented querying
  - Imperative code
- SQL is declarative:
  - Programmer explains a solution
  - System figures out the best way to find the solution
- Use declarative query languages for document databases

### Query Languages

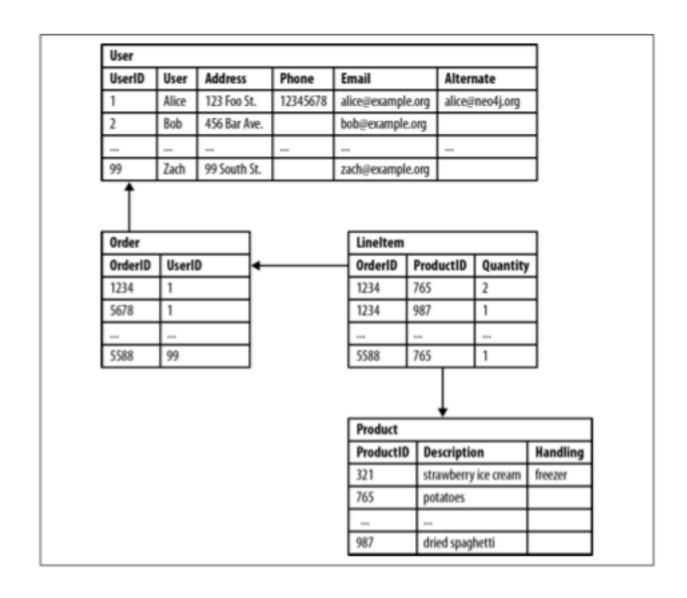
- Map-Reduce (neither declarative nor imperative):
  - Consists of only two pieces of code
    - Mapping: Selecting from Documents
    - Reducing: Take selection elements and operate on them

# Alternatives to Relational Schemes: Graph Models

- Graphs consists of vertices and edges
  - Example:
    - Social graphs: vertices are people and edges are relationships such "knows"
    - Web graph: vertices are pages and edges are links
    - Road networks: vertices are places and edges are connections

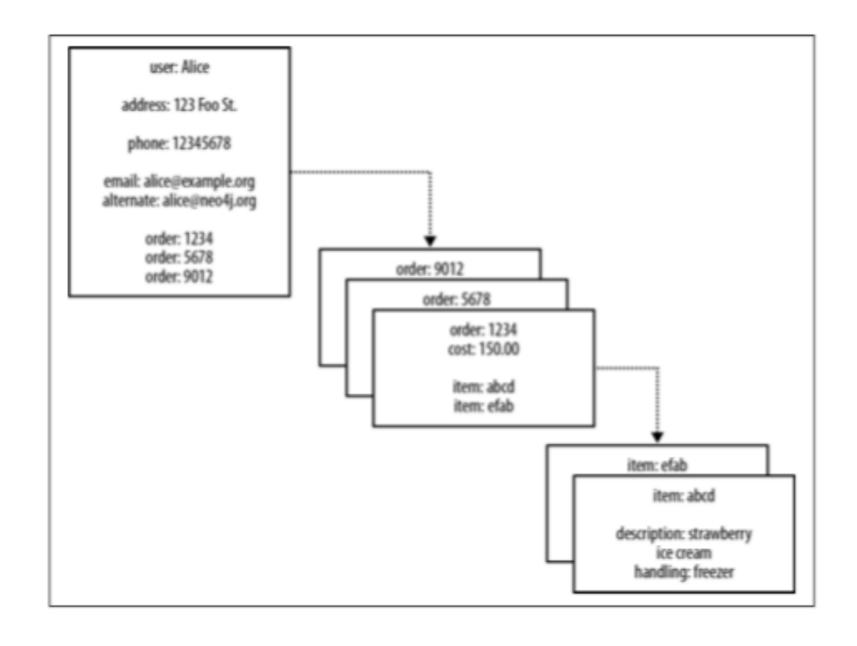
# Alternatives to Relational Schemes: Graph Models

Relational Database hides semantic relationships

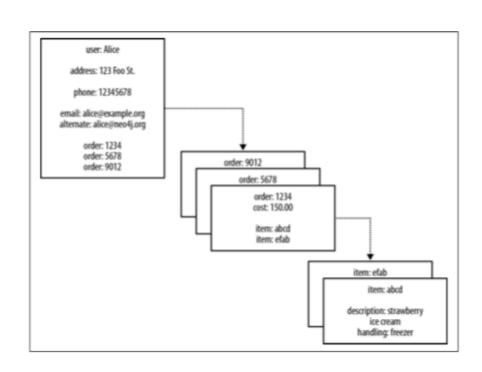


# Alternatives to Relational Schemes: Graph Models

Document model hides semantic relationships

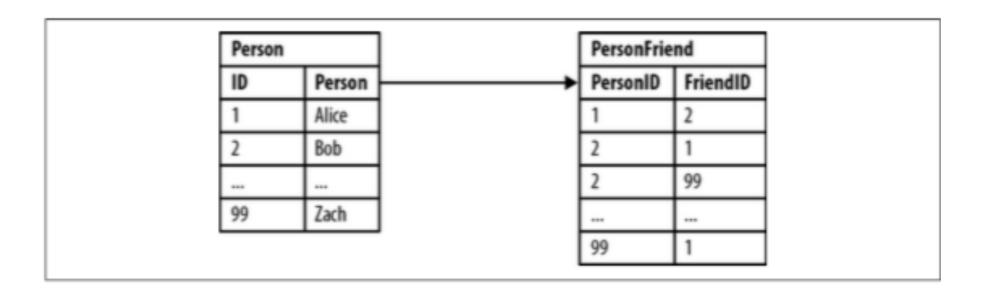


- Some property values are really references to foreign aggregates
  - Aggregate's identifier is a foreign key
- Relationships between them are not explicitly accessible
  - Joining aggregates becomes expensive



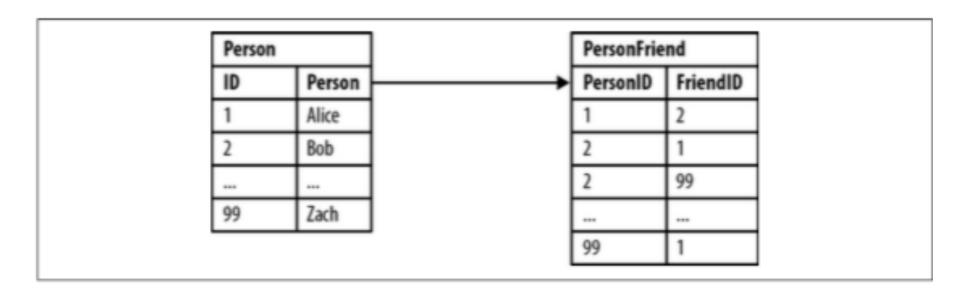
- Relational Database
  - Some queries are simple:

```
SELECT pl.Person
FROM Person pl JOIN PersonFriend
ON PersonFriend.FriendID = pl.ID JOIN Person p2
ON PersonFriend.PersonID = p2.ID WHERE p2.Person = 'Bob'
```



- Relational Database
  - Some queries others are more involved: Friends of Bob

```
SELECT p1.Person
FROM Person p1 JOIN PersonFriend
   ON PersonFriend.PersonID = p1.ID JOIN Person p2
   ON PersonFriend.FriendID = p2.ID
WHERE p2.Person = 'Bob'
```



- Relational Database
- Some queries others are difficult: Alice's friends of friends
  SELECT pl.Person AS PERSON, pl.Person AS FRIEND\_OF\_FRIEND FROM

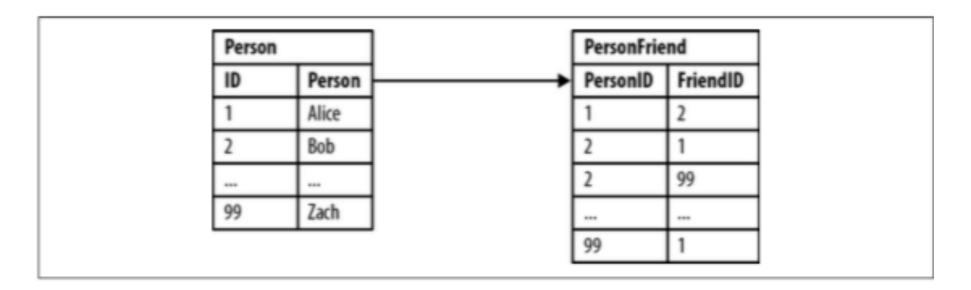
```
PersonFriend pf1 JOIN Person p1
```

```
ON pf1.PersonID = p1.ID JOIN PersonFriend pf2
```

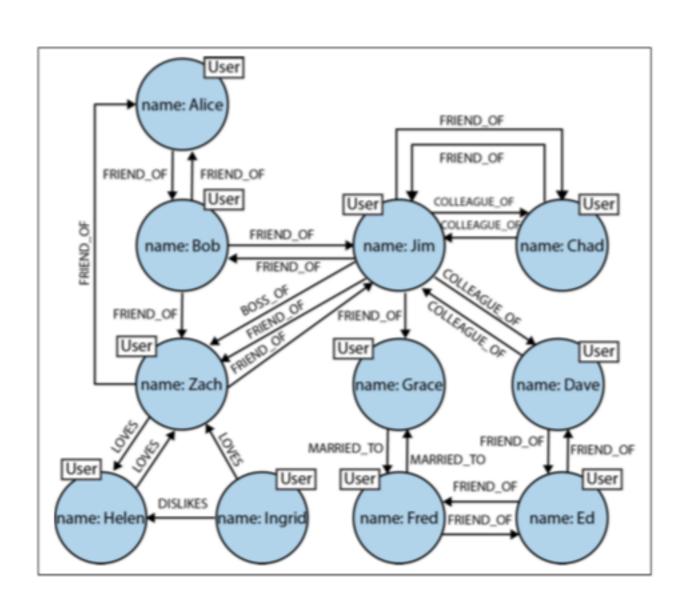
ON pf2.PersonID = pf1.FriendID JOIN Person p2

ON pf2.FriendID = p2.ID

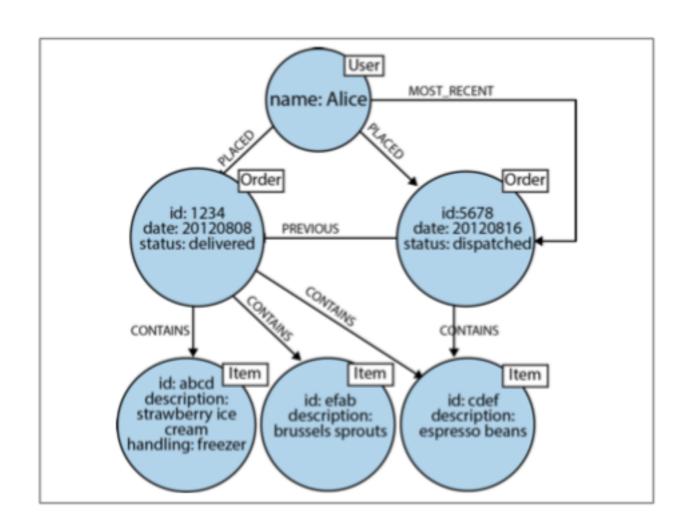
WHERE pl.Person = 'Alice' AND pf2.FriendID <> pl.ID



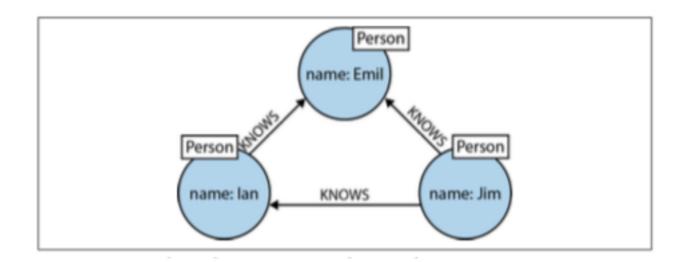
- Property graph model by Neon
  - Each vertex consists of
    - A unique identifier
    - A set of outgoing edges
    - A set of incoming edges
    - A collection of properties key-value pairs
  - Each edge consists of
    - A unique identifier
    - The tail vertex
    - The head vertex
    - A label to describe the relationship
    - A collection of properties key-value pairs



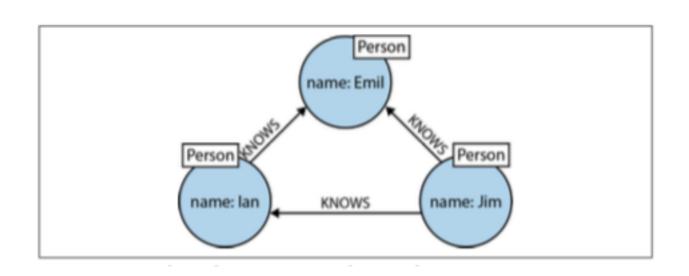
Order history as a property graph



- Processing queries in Neo4j
  - Use Cypher (from "The matrix")
  - Can describe a path



(emil) <-[:KNOWS] - (jim) - [:KNOWS] -> (ian) - [:KNOWS] -> (emil)



```
(emil:Person {name:'Emil'})
    <-[:KNOWS]-(jim:Person {name:'Jim'})
    -[:KNOWS]->(ian:Person {name:'Ian'})
    -[:KNOWS]->(emil)
```

Finding the mutual friends of Jim:

```
MATCH (a:Person {name:'Jim'})-[:KNOWS]->(b)-[:KNOWS]->(c), (a)- [:KNOWS]->(c)
RETURN b, c
```

Triple Stores

- Information is stored as (subject, predicate, object)
  - Subjects correspond to vertices
  - Objects are
    - A value in a primitive data type (jim : age : 64)
    - Another vertex (jim : friend\_of : thomas)

```
@prefix : </example>
:lucy a
                    :Person
:lucy :name
                   "Lucy"
:lucy :born in
                   :idaho
:idaho a
                   :Location
:idaho :name
                   "Idaho"
:idaho :type
                   "State"
        :within
:idaho
                   :usa
```

- Triple stores are the language of the semantic web
- Semantic web:
  - Machine readable description of type of links
    - e.g. image, text, ...
  - Creates web of data a database of everything
- Stored in Resource Description Framework (RDF)
- SPARQL query language for triple stores